

# Discrete search in design optimization

Martin Fuchs

CERFACS, Toulouse, France

in collaboration with Arnold Neumaier, University of Vienna,  
Austria

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## Design optimization

- black box cost function
  - simulation based
  - computationally expensive evaluation
  - strong nonlinearities, discontinuities
  - hidden constraints
- various design choices
  - continuous, discrete, and categorical choices
  - mixed variables



## Problem formulation

$$\begin{aligned} \min_{\theta, z} \quad & F(z) \\ \text{s.t.} \quad & z = Z(\theta), \\ & \theta \in \mathbf{T}. \end{aligned}$$

- $F : \mathbb{R}^{n_z} \rightarrow \mathbb{R}$  black box
- $\mathbf{T}$  is the domain of possible design choices  $\theta$
- $Z$  is the design selection mapping



## Problem generalization

$$\begin{aligned} \min_{\theta, x} \quad & c^T x \\ \text{s.t.} \quad & F(Z(\theta)) \leq Ax, \\ & \theta \in \mathbf{T}. \end{aligned}$$

- special case  $c = A = 1$  gives former formulation
- black box input  $z$  substituted by  $z = Z(\theta)$



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## Design selection example: discrete case

$\theta$	Thruster	$F/N$	$I_{sp}/s$	$m_{thrust}/kg$
1	Aerojet MR-111C	0.27	210.0	200
2	EADS CHT 0.5	0.50	227.3	200
3	MBB Erno CHT 0.5	0.75	227.0	190
4	TRW MRE 0.1	0.80	216.0	500
5	Kaiser-Marquardt KMHS Model 10	1.0	226.0	330

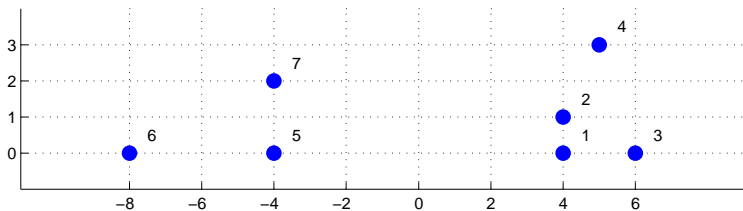
- contains specifications of design components and the associated choice variable  $\theta$
- the table mapping  $Z : \theta \rightarrow (F, I_{sp}, m_{thrust})$  assigns an input parameter vector to a given design point  $\theta$
- in general  $F$  is a physical model based on  $Z(\mathbf{T})$ , rather than on  $\mathbf{T}$





## Search space $\mathbf{T}$

- $\mathbf{T}$  is the set of all possible designs
- $\mathbf{T} = T^1 \times T^2 \times \dots \times T^{n_0}$
- $T^i = \begin{cases} \{1, 2, \dots, N_i\} & \text{in the discrete case,} \\ [\underline{\theta}^i, \overline{\theta}^i] & \text{in the continuous case.} \end{cases}$
- for discrete  $T^i$  the design selection mapping provides a finite multidimensional set  $Z^i(T^i)$ , e.g.,



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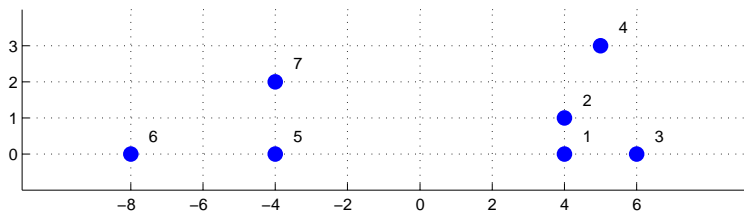
## Heuristic approaches

- SNOBFIT (fits a quadratic model of the objective function and minimizes this model)
- Evolutionary algorithms
- Separable underestimation
- **Splitting based on convex relaxation**
- combination with methods for continuous variables



## Convex relaxation based splitting: idea

- remember:  $F$  is based on  $Z(\mathbf{T})$  rather than on  $\mathbf{T}$



- for discrete  $T^i$  relaxation to the convex hull of  $Z^i(T^i)$



## Convex relaxation of $Z(\mathbf{T})$

$$\min_{z,v,\lambda} c^T x$$

$$\text{s.t. } F(z) \leq Ax,$$

$$z = (v^1, \dots, v^{n_0}),$$

$$v^i = \sum_{j=1}^{N_i} \lambda_j^i Z^i(j) \text{ for } i \in I_d,$$

$$\sum_{j=1}^{N_i} \lambda_j^i = 1 \text{ for } i \in I_d,$$

$$\lambda_j^i \geq 0 \text{ for } i \in I_d, 1 \leq j \leq N_i,$$

$$v^i \in [\underline{\theta}^i, \overline{\theta}^i] \text{ for } i \in I_c.$$

} convex  
combination



## Approximate linear solution

$$\min_{z, x, \mu, v, \lambda} c^T x + \varepsilon \|\mu\|_p$$

$$\text{s.t.} \quad \sum_{j=1}^{N_0} \mu_j F_j \leq Ax,$$

$$z = \sum_{j=1}^{N_0} \mu_j z_j,$$

$$\sum_{j=1}^{N_0} \mu_j = 1,$$

$$z = (v^1, \dots, v^{n_0}),$$

$$v^i = \sum_{j=1}^{N_i} \lambda_j^i Z^i(j) \text{ for } i \in I_d,$$

$$\sum_{j=1}^{N_i} \lambda_j^i = 1 \text{ for } i \in I_d,$$

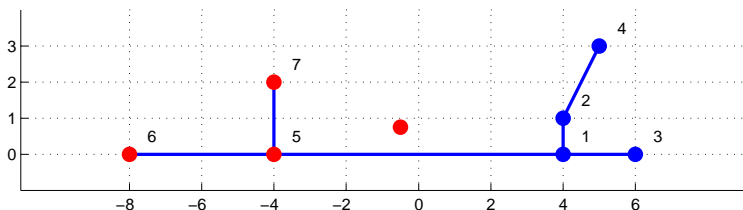
$$\lambda_j^i \geq 0 \text{ for } i \in I_d, 1 \leq j \leq N_i,$$

$$v^i \in [\underline{\theta}^i, \overline{\theta}^i] \text{ for } i \in I_c.$$

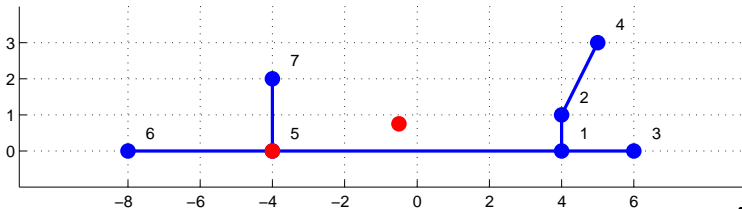
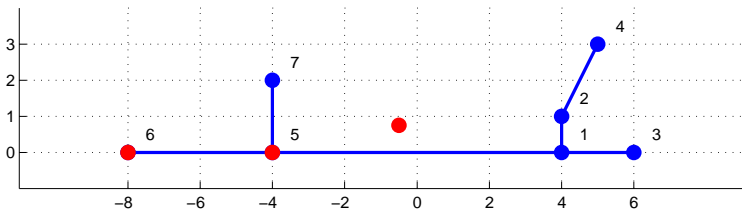


# Splitting

- use the coefficients of the convex relaxation as weights on the minimum spanning tree (MST) of  $Z(\mathbf{T})$
- split the MST in two of parts of similar total weight
- Example:



## Splitting ctd.





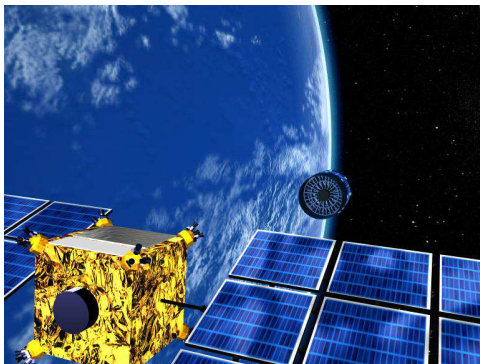
## Solver strategy

- ① Find a relaxed approximate solution on the current branch.
- ② Round the relaxed solution  $\hat{z}$  to the next feasible point, i.e.,  $\hat{z}_{\text{round}} := \arg \min_{\{z \in Z(\mathbf{T})\}} \|z - \hat{z}\|_2$ .
- ③ Start neighborhood search from  $\hat{z}_{\text{round}}$ .
- ④ Split on the variable with maximal deviation during Step 3.
- ⑤ Select the branch with the best function value in Step 3 for the next iteration.

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## XEUS mission – permanent space-borne X-ray observatory



- complex design problem, 10 dimensions
- $4 \times 14 \times 6 \times 8 \times 5 \times 20 \times 9 \times 44 \times 30 \geq 3 \cdot 10^9$   
discrete choices
- 1 continuous choice variable



## Optimization results

- total mass  $m = 1566$  kg
- found in 4 out of 5 runs of 2500 function evaluations each
- 1 run failed because we found no feasible starting point
- previous study used  $\geq 50000$  function calls



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## Summary

- Exploit structural knowledge about the discrete search space.
- Speed up the splitting procedure in branching algorithms.
- Solve successfully higher dimensional real-life design optimization problems.

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